

## Logic Programming Languages

- *Logic* programming systems allow the programmer to state a collection of axioms from which theorems can be proven.
- Express programs in a form of symbolic logic
- Use a logical inferencing process to produce results
- *Logic* programming languages are rule based languages
- *Declarative* rather than *procedural*:
  - Only specification of *results* are stated (not detailed *procedures* for producing them)

## Introduction to predicate calculus

- Particular form of symbolic logic used for logic programming is **predicate calculus**. Symbolic logic provides basis for logic programming. predicate calculus contains Propositions

### Propositions

- **Proposition** is logical statement that may or may not be true. It consists of objects and relationships of objects to each other

### Objects

- Objects in propositions are represented by simple terms: either constants or variables
- *Constant*: a symbol that represents an object
- *Variable*: a symbol that can represent different objects at different times
  - Different from variables in imperative languages
- The simplest form of proposition is called **atomic proposition** consist of compound terms. A **compound term** is one element of a mathematical relation, written in a form that has the appearance of mathematical function

Compound term composed of two parts

- Functor: function symbol that names the relationship
- Ordered list of parameters (tuple)
- Examples:
  - student(jon)      --- jon is student
  - like(rony, OS)
  - like(nick, windows)    ---- nick like windows
  - like(jim, linux)

Propositions are connected by operators

Name	Symbol	Example	Meaning
Negation	$\neg$	$\neg a$	not a
Conjunction	$\wedge$	$a \wedge b$	a and b
Disjunction	$\vee$	$a \vee b$	a or b
Equivalence	$\equiv$	$a \equiv b$	a is equivalent to b
Implication	$\supset$ $\subset$	$a \supset b$ $a \subset b$	a implies b b implies a

## Horn clauses

Propositions can be stated as Horn clauses.

- In almost all logic languages, propositions are written in a standard form known as Horn clause and it consists of head (H) and body ( $B_1, \dots, B_n$ )

following is the general form of horn clause

$H :- B_1, \dots, B_n.$  read as  $B_1, \dots, B_n$  implies  $H$

When  $B_i$  are all true, we can deduce that  $H$  is true *i.e Antecedent implies Consequent.* ( $B_1, \dots, B_n$  implies  $H$ )

*Here comma ( , ) is taken as AND( $\wedge$ ) operator*

- *Antecedent*: right side ( $B_1, \dots, B_n$ )
- *Consequent*: left side ( $H$ )
- Antecedent is called as if part and Consequent is called as then part*

- Horn clauses Basis for logic programming

- **Horn clause - can have only two forms**

- **Headed Horn clause**: single atomic proposition on left side

Ex:  $\text{parent}(X,Y):- \text{mother}(X,Y).$

If X is a mother of Y then X is a parent of Y

$\text{grandparent}(X,Z):- \text{parent}(X,Y), \text{parent}(Y,Z)$

If X is a parent of Y and Y is a parent of Z then X is a grandparent of Z

- **Headless Horn clause**: empty left side (used to state facts)

Ex:  $\text{mammal}(\text{human})$  ----- human is a mammal  
 $\text{female}(\text{uma}).$

- **Resolution** an inference rule for horn clauses that allows inferred propositions to be computed from given propositions
- Resolution is the primary activity of a Prolog interpreter
- In order to derive new statements, logic programming system combines existing statements, cancelling like terms, through a process known as resolution

Ex:  $C :- A, B$

$D :- C$

---

$D :- A, B$

**Unification** is a process of finding values for variables in propositions that allows matching process to succeed. *Unification make two terms identical*

- **Instantiation**: assigning temporary values to variables to allow unification to succeed
- After instantiating a variable with a value, if matching fails, may need to *backtrack* and instantiate with a different value

- Ex: **Prolog unifies  $\text{man}(X)$  with  $\text{man}(\text{fred})$** , thereby instantiating the variable X to fred .

- *Unification* make two terms (*X* and *fred*) identical
- After instantiating a variable with a value, if matching fails, may need to *backtrack* and instantiate with a different value

## Programming with prolog

**Prolog** is a logic programming language and it is developed at University of Aix-Marseille and University of Edinburgh

### The basic elements of prolog

**Terms**

**Fact statements**

**Rule Statements**

**Goal Statements**

**Simple Arithmetic**

**List structures**

### Terms

All Prolog statements, as well as Prolog data, are constructed from **terms**.

**Term** is a constant, variable, or structure

- *Constant* is either an atom or an integer ex: fred, jake, 23 etc
- *Atom* is a symbolic value of Prolog
- Atom consists of either:
  - a string of letters, digits, and underscores beginning with a lowercase letter
  - a string of printable ASCII characters delimited by apostrophes
- *Variable*: any string of letters, digits, and underscores beginning with an uppercase letter  
Ex: X , Y , Z etc
- Structures represent the atomic propositions of predicate calculus, and their general form is functor (parameter list) ex: parent (X,Y) where X and Y are term1 and term2
- The functor is any atom and is used to identify the structure. The parameter list can be any list of atoms, variables,

### Fact statements

Fact statements are used for the hypotheses

Headless Horn clauses are called Fact statements

The simplest form of headless Horn clause in Prolog is a single structure, which is interpreted as an unconditional assertion, or fact.

## Examples of fact statements

```
mammal(human)
female(shelley).
male(bill).
father(bill, jake).
mother(suma, shelley).
```

## Rule Statements

- Used for the hypotheses
- Headed Horn clause are called as Rule Statements
- The general form of the Prolog headed Horn clause statement is

consequence :- antecedent\_expression.

It is read as follows: "consequence can be concluded if the antecedent expression is true"

- Right side: *antecedent* (**if** part)
  - May be single term or conjunction
- Left side: *consequent* (**then** part)
  - Must be single term

*Conjunction*: multiple terms separated by logical AND operations

## Examples of Rule Statements

```
ancestor(mary,shelley):- mother(mary,shelley).
parent(X,Y):- mother(X,Y).
parent(X,Y):- father(X,Y).
grandparent(X,Z):- parent(X,Y), parent(Y,Z).
if X is a parent of Y AND Y is a parent of Z then X is a grandparent of Z
sibling(X,Y):- mother(M,X), mother(M,Y),
               father(F,X), father(F,Y).
```

## Goal Statements

- For theorem proving, theorem is in form of proposition that we want system to prove or disprove. These propositions are called *goal statements* or **queries**
- Goal Statements are Same format as headless Horn clauses

For example, we could have

```
man(fred)
```

to which the system will respond either `yes` or `no`. The answer `yes` means that the system has proved the goal was true under the given database of facts and relationships.

- Conjunctive propositions and propositions with variables also legal goals  
`father(X,mike)`

The approach of logic programming is to use as a database a collection of facts and rules that state relationships between facts and to use an automatic inferencing process to check the validity of new propositions, assuming the facts and rules of the database are true. This approach is the one developed for automatic theorem proving.

## Simple Arithmetic

- Prolog supports integer variables and integer arithmetic
- `is` operator: takes an arithmetic expression as right operand and variable as left operand  
`A is B / 17 + C`

## Trace

- Built-in structure that displays instantiations at each step
- *Tracing model* of execution - four events:
  - *Call* (beginning of attempt to satisfy goal)
  - *Exit* (when a goal has been satisfied)
  - *Redo* (when backtrack occurs)
  - *Fail* (when goal fails)

## Inferencing Process of Prolog

- Queries are called goals
- To prove a goal is true, must find a chain of inference rules and/or facts. For goal Q:  
`B :- A`  
`C :- B`  
...  
`Q :- P`
- The following example illustrates the difference between forward and backward chaining. Consider the query:  
`man(bob) .`  
Assume the database contains  
  
`father(bob) .`  
`man(X) :- father(X) .`
- **Forward chaining** would search for and find the first proposition. The goal is then inferred by matching the first proposition with the right side of the second rule (`father(X)`) through instantiation of `X` to `bob` and then matching the left side of the second proposition to the goal.
- **Backward chaining** would first match the goal with the left side of the second proposition (`man(X)`) through the instantiation of `X` to `bob`. As its last step, it would match the right side of the second proposition (now `father(bob)`) with the first proposition

## List Structures (lists)

- Other basic data structure (besides atomic propositions we have already seen): list
- *List* is a sequence of any number of elements
- Elements can be atoms, atomic propositions, or other terms (including other lists)

[apple, prune, grape, kumquat]

[] (*empty list*)

[H | T] (*head H and tail T*)

if [ H | T ] = [ 1, 2,3], then H = 1 and T = [ 2,3]

## Execution in Prolog

database contains facts and rules

Following is database

likes(mary,food).

likes(mary,wine).

likes(john,wine).

- The following queries yield the specified answers.

| ?- likes(mary,food).

the system will respond as

O/P : yes.

| ?- likes(john,wine).

yes.

| ?- likes(john,food).

no.

## Applications of Logic Programming

- Relational database management systems Relational database management systems (RDBMSs) store data in the form of tables. Queries on such databases are often stated in Structured QOquery Language (SQL). SQL is nonprocedural in the same sense that logic programming is nonprocedural. The user does not describe how to retrieve the answer; rather, he or she describes only the characteristics of the answer. The connection between logic programming and RDBMSs should be obvious. Simple tables of information can be described by Prolog structures, and relationships between tables can be conveniently and easily described by Prolog rules. The retrieval process is inherent in the resolution operation. The goal statements of Prolog provide the queries for the RDBMS. Logic programming is thus a natural match to the needs of implementing an RDBMS.

**One of the advantages of using logic programming** to implement an RDBMS is that only a single language is required. In a typical RDBMS, a database language includes statements for data definitions, data manipulation, and queries. The general-purpose language is used for processing the data and input and output functions. All of these functions can be done in a logic programming language.

**Another advantage of using logic programming** to implement an RDBMS is that deductive capability is built in. Conventional RDBMSs cannot deduce anything from a database other than what is explicitly stored in them. They contain only facts, rather than facts *and* inference rules. The primary disadvantage of using logic programming for an RDBMS, compared with a conventional RDBMS, is that the logic programming implementation is slower. Logical inferences simply take longer than ordinary table look-up methods using imperative programming techniques.

- **16.8.2 Expert Systems**

Expert systems are computer systems designed to imitate human expertise in some particular domain. They consist of a database of facts, an inferencing process, some heuristics about the domain, and some friendly human interface that makes the system appear much like an expert human consultant. In addition to their initial knowledge base, which is provided by a human expert, expert systems learn from the process of being used, so their databases must be capable of growing dynamically. Also, an expert system should include the capability of interrogating the user to get additional information when it determines that such information is needed.

One of the central problems for the designer of an expert system is dealing with the inevitable inconsistencies and incompleteness of the database. Logic programming appears to be well suited to deal with these problems. For example, default inference rules can help deal with the problem of incompleteness.

Prolog can and has been used to construct expert systems. It can easily fulfill the basic needs of expert systems, using resolution as the basis for query processing, using its ability to add facts and rules to provide the learning capability, and using its trace facility to inform the user of the “reasoning” behind a given result. Missing from Prolog is the automatic ability of the system to query the user for additional information when it is needed.

One of the most widely known uses of logic programming in expert systems is the expert system construction system known as APES. The APES system includes a very flexible facility for gathering information from the user during expert system construction. It also includes a second interpreter for producing explanations to its answers to queries.

APES has been successfully used to produce several expert systems, including one for the rules of a government social benefits program and one for the British Nationality Act, which is the definitive source for rules of British citizenship.

### **16.8.3 Natural-Language Processing**

Certain kinds of natural-language processing can be done with logic programming. In particular, natural-language interfaces to computer software systems, such as intelligent databases and other intelligent knowledge-based systems, can be conveniently done with logic programming. For describing language syntax, forms of logic programming have been found to be equivalent to context-free grammars. Proof procedures in logic programming systems have been found to be equivalent to certain parsing strategies. In fact, backward-chaining resolution can be used directly to parse sentences whose structures are described by context-free grammars. It has also been discovered that some kinds of semantics of natural languages can be made clear by modeling the languages with logic programming

# Multiparadigm Languages

A new class of programming languages and environments is being developed to help solve this problem. They do not restrict the programmer to only one paradigm (see box on p. 8); incorporating two or more of the conventional program paradigms. For example, the Loops system, described in the first article in this issue, combines features of the Lisp, functional, rule oriented, and object-oriented paradigms. These multiparadigm systems are being created to give the programmer the right tool at the right time.

The most common multiparadigm system is the conventional operating system, which embodies several different programming languages.

## Some programming paradigms

**Access oriented:** The specification of side-effects or demons attached to the manipulation of variables [an extension of Loops].

**Data-structure-oriented:** Approaching a solution by way of a single powerful data structure, such as lists (Lisp), arrays (APL), sets (SETL), relational databases (SQL).

**Functional:** A pure mathematical specification of the solution to a problem, eliminating the conventional von Neumann model of memory and variables [pure Lisp, Backus's FP language].

**Imperative:** Sequential, block-structured commands with static scoping of variables [Fortran, Pascal, C].

**Object-oriented:** Grouping data into objects or abstract data types, where each object (or class of objects) has a set of operations (methods) to manipulate the data stored in that object [Smalltalk, Simula, CLU].

**Parallel:** Specifying multiple processes in the context of one processor or a distributed collection of processors [Modula-2, Ada, NIL, CSP].

**Real-time:** Specifying the constraints associated with physical objects such as robot arms, missiles, and analog devices [Arctic, AML, an extension to Lucid].

**Rules-oriented:** Specifying the constraints of the problem, rather than the algorithm for finding a solution [Prolog, OPS5].

The languages in brackets exhibit constructs that belong to the indicated paradigm. Some of the languages, especially the fourth-generation languages, are multiparadigm in that they incorporate features from different categories.



## **Building a multiparadigm system**

There are at least four ways to build a multiparadigm language. The first is to laminate together existing languages, for example, combining the syntax (and semantics) of Prolog, Lisp, and C into one system. An advantage of such a system is that users are already familiar with at least one component of the system, so they can start using the system quickly. A serious disadvantage is the complex semantics caused by unintended interactions between the component language constructs.

